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## **Accelerated Hard Thresholding Algorithms for Sparse Recovery**

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**Abstract.** In linear observations, i.e., a system of linear equations Ax = b ( $A \in \mathbb{R}^{M \times N}$ ), the hard thresholding pursuit (HTP) is used to find a sparse signal. HTP is an iterative algorithm that has been found many applications in compressive sensing, due to its good recovery performance, which includes linear convergence speed, high recovery rate, and stability. In this paper, we further develop accelerated algorithms to deal with a inear least square (LLS) problem in each iteration. Theoretically, we prove that all these algorithms are convergent, provided that the sensing matrix has suitable restricted isometry property. Numerical experiments on sparse signal recovery demonstrate the efficiency of the proposed methods.

## Introduction

The compressive sensing [1,2] aims to reconstruct a sparse signal  $x \in \mathbb{R}^N$  from a few linear measurements b = Ax ( $A \in \mathbb{R}^{M \times N}$ ). Such problem frequently appeared in signal and image processing, face recognition, and more [3,4]. Many literatures have made great contributions to this research [1,2,3,4]. As is well known, the basis pursuit (BP) [1], which is one of the most famous reconstruction methods in , is to solve the following  $l_1$ -minimization problem,

$$\min_{x \in \mathcal{R}^N} \parallel x \parallel_1 st. Ax = y, \tag{1}$$

where  $A \in \mathcal{R}^{M \times N}$  is the sensing matrix and  $\mathcal{Y} \in \mathcal{R}^{M}$ . Problem (1) is a convex problem, and then many convex methods are available. Although many methods are developed, solving BP still requires a lot of time. Besides the basic pursuit, paper [5] proposed an alternative strategy based on learning the sparsity of the signal, called hard thresholding pursuit (HTP), which is quite fast (converges linearly) and stable. The hard thresholding pursuit, which combines the hard thresholding strategy of iterative hard thresholding algorithm (IHT) [6] and the projection method of compressive sampling matching pursuit algorithm (CoSaMP) [7], can be described as follows:

$$S_{n+1} = C_s(x_n + A^*(y - Ax_n)), HTP_1$$

$$x_{n+1} = arg\min_{z} \| y - Az \|_2, supp(z) \subseteq S_{n+1}, HTP_2$$

where  $x_n$  is the  $n^{th}$  iterate and  $C_s(x)$  denotes the indices of s largest absolute entries of x.

Here, we introduce the definition of Restricted Isometry Property (RIP) which plays an important role in analyzing the convergence of HTP algorithm. The  $^S$  order restricted isometry constant  $\delta_s = \delta_s(A)$  of  $A \in \mathcal{R}^{M \times N}$  has been defined in [8] as the smallest  $\delta$  such that

$$(1 - \delta) \| x \|_2^2 \le \| Ax \|_2^2 \le (1 + \delta) \| x \|_2^2, \tag{2}$$

for any S-sparse vector  $x \in \mathbb{R}^N$ . The author in [5] has proved that the HTP algorithm has a linear convergent rate provided  $\delta_{3s} < \frac{\sqrt{3}}{3}$ .

In the second step of each iteration of HTP algorithm (HTP 2), a linear least square (LLS) problem needs to be solved. However, as the scale of the problem increases, solving LLS directly will be very



expensive sometimes even intractable. Due to this, HTP algorithm will encounter challenge if the scale of the problem and the sparsity level are all very large. A scheme called fast hard thresholding pursuit (FHTP) proposed in [5] provides a method in large scale case. The main idea of FHTP is generating an approximate solution by applying the gradient method to the LLS problem in HTP 2 rather than solving it exactly. Here, we present the detailed procedure of this scheme:

Algorithm 1 Fast Hard Thresholding Pursuit<sup>5</sup>

```
 \begin{aligned} & \textbf{Require: parameters } h_{n+1}^i, k \\ & \textbf{Initialization: } y^0 = 0 \\ & \textbf{for } n = 0, 1, 2, \dots \\ & S_{n+1} = C_s(x_n + A^*(y - Ax_n)) \\ & x_{n+1}^0 = (x_n + A^*(y - Ax_n))_{S_{n+1}} \\ & y_{n+1}^1 \in \mathcal{R}^s, y_{n+1}^1 = (x_{n+1}^0)_{S_{n+1}} \\ & \textbf{for } i = 1, 2, \dots, k-1 \\ & A_{n+1} \in \mathcal{R}^{n \times s}, A_{n+1} = A(:, S_{n+1}) \\ & y_{n+1}^{i+1} = y_{n+1}^i + h_{n+1}^i A_{n+1}^*(y - A_{n+1}y_{n+1}^i) \\ & \textbf{end for} \\ & x_{n+1} \in \mathcal{R}^n, (x_{n+1})_{S_{n+1}} = y_{n+1}^k, (x_{n+1})_{\overline{S_{n+1}}} = 0 \\ & \textbf{end for} \end{aligned}
```

where  $h_{n+1}^i$  is the step parameter. The FHTP also has a linear convergence rate if the iterative number

k in the loop of Algorithm 1 is big enough and  $\delta_{3s} < \frac{\sqrt{3}}{3}$ . Paper [5] proves that when  $h_{n+1}^i \equiv 1$  the following inequality holds,

$$\| x_{n+1} - x \|_{2} \le \sqrt{\frac{2\delta_{3s}^{2} + \delta_{3s}^{2k+2}(1 - 3\delta_{3s}^{2})}{1 - \delta_{3s}^{2}}} \| x_{n} - x \|_{2}, \tag{3}$$

where x is any S-sparse and obeys Ax = b,  $\{x_n\}_{n=0,1,2,\dots}$  are obtained by FHTP.

Note that HTP 2 is actually a optimization problem, which can be solved by faster optimization method such as Nesterov method [9, 10], from the perspective of optimization. While, HTP 2 is also to solve a positive definition linear system, which can be solved by fast linear algebra method such as conjugate gradient method [11], from the perspective of linear algebra. In this paper, we show that if the gradient method applied in FHTP 2 is replaced by other methods, the scheme is also convergent. Therefore, we can employ some other fast methods to get faster HTP schemes. The rest of the paper is organized as follows. Section 2 introduces the preliminary knowledge. Section 3 presents the algorithms and their convergence. Section 4 contains numerical results. Finally, Section 5 concludes the paper.

#### **Notations**

In this paper, 
$$A = \{A_1, A_2, ..., A_N\} \in \mathcal{R}^{M \times N}$$
,  $x = \{x(1), x(2), ..., x(n)\}^* \in \mathcal{R}^N$  and  $S = \{i_1, i_2, ..., i_s\} \subseteq \{1, 2, ..., N\}$ .  $\overline{S}$  denotes as the set  $\{1, 2, ..., N\} \setminus S$ ,  $x_S \in \mathcal{R}^N$  denotes as  $x_S(i) = h(i), i \in S, x_S(i) = 0, i \in \overline{S}$ .  $A(:, S) \in \mathcal{R}^{M \times k}$  denotes as  $A(:, S) = \{A_{i_1}, A_{i_2}, ..., A_{i_s}\}$ .

#### **Preliminaries**

In this part, we introduce two powerful algorithms which have been widely used in optimization and numerical linear algebra. These two algorithms will be employed in the proposed algorithm to accelerate the speed of HTP.

**2.1 The Nesterov method.** The Nesterov method can be used to accelerate the speed of gradient method in minimizing the smooth convex function problem

$$\min f(x),$$
 (4)

where f has a Lipschitz continuous gradient with constant L = L(f), i.e.,  $\|\nabla f(x_1) - \nabla f(x_2)\|_2 \le L \|x_1 - x_2\|_2, x_1, x_2 \in dom(f).$  (5)



There are many variants of the Nesterov method [12, 13, 14]. Here, we present a simple one appeared in [14] as follows:

## Algorithm 2 The Nesterov method<sup>14</sup>

```
Require: parameters h > 0

Initialization: y^0 = x^0, \theta_0 = 1

for k = 0, 1, 2, ...

x^{k+1} = y^k - h_k \nabla f(y^k)

\beta_{k+1} = \frac{1}{2}(1 - \theta_k)(\sqrt{\theta_k^2 + 4} - \theta_k)

y^{k+1} = x^k + \beta_k(x^{k+1} - x^k)

\theta_{k+1} = \frac{1}{2}\theta_k(\sqrt{\theta_k^2 + 4} - \theta_k)
```

Actually, this algorithm is equivalent to Constant Step Scheme II on Page 80 of [9] and FISTA on Page 193 of [12] without the nonsmooth regularization function. The  $h_k$  is the step parameter, a usually choice is  $h_k = \frac{1}{L}$ . The author in [9] shows that

usually choice is 
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. The author in [9] shows that 
$$f(y^k) - f^* \le \frac{2L\|y^0 - y^*\|_2^2}{(k+1)^2},$$
 (6)

where  $y^*$  is a solution to (4) and  $f^* = f(y^*)$ . When apply Algorithm 2 to  $HTP_2$ , the objective function is  $f(x) = ||y - A_n x||_{2, \text{ where }}^2 A_n = A(:, S_n) \in \mathbb{R}^{M \times s}$ . Because A has the restricted isometry property, it is easy to obtain that

$$\|A_n x\|_2^2 \le (1 + \delta_s) \|x\|_2^2 < 2, \forall x \in \mathcal{R}^s.$$
 (7)

That means  $\|A_n^*A_n\|_2 < 2$ , which also means that

$$\|\nabla f(x_1) - \nabla f(x_2)\|_2 = \|A_n^* A_n(x_1 - x_2)\|_2 < 2 \|x_1 - x_2\|_2.$$
(8)

Then, we can set  $h_k = \frac{1}{2}$  in Algorithm 4 (see below).

# **2.2 Conjugate Gradient Method.** The HTP 2 is essentially solving following linear system $A_n^*A_nx = A_n^*y$ , (9)

where  $A_n = A(:, S_n) \in \mathbb{R}^{M \times s}$  and  $x \in \mathbb{R}^s$ . Because A has the restricted isometry property, we can immediately get two conclusions: the first is that  $A_n^* A_n$  is positive definition; the second is that  $cond(A_n^* A_n) \leq \frac{1+\delta_s}{s}$ 

 $cond(A_n^*A_n) \le \frac{1+\delta_s}{1-\delta_s}$ . Based on these two points, we can apply conjugate gradient method (CG) whose convergence rate closely depends on the condition number of the coefficient matrix to (9).

We emphasize that the matrix-matrix-vector multiplication  $A_n^* A_n x$  should be computed as following scheme:  $\overline{Z} \to A_n x$ ,  $Z \to A_n^* \overline{Z}$ . A linear convergent rate about CG method has been proposed by [15] as

$$\parallel x_k - x^* \parallel_{A_n^* A_n} \le 2 \left( \frac{\sqrt{\kappa} - 1}{\sqrt{\kappa} + 1} \right)^k \parallel x_0 - x^* \parallel_{A_n^* A_n}, \tag{10}$$

Algorithm 3 Conjugate gradient method to (9)<sup>15</sup>

```
Require: Initialization x_0 \in \mathcal{R}^s, compute r_0 = A_n^*y - A_n^*A_nx_0, let p_0 = r_0 for k = 0, 1, 2, \dots compute \alpha_k = \frac{\langle r_k, r_k \rangle}{\langle A_n p_k, A_n p_k \rangle} x_{k+1} = x_k + \alpha_k p_k r_{k+1} = r_k - \alpha_k A_n^*A_n p_k \beta_k = \frac{\langle r_{k+1}, r_{k+1} \rangle}{\langle r_k, r_k \rangle} p_{k+1} = r_{k+1} + \beta_k p_k end for
```

where  $x^*$  is the exact solution to (9),  $\|x\|_{A_n^*A_n} = x^T A_n^* A_n x$  and  $\kappa$  is the condition number of  $A_n^*A_n$ . Considering  $A_n$  is a submatrix of A which is restricted isometry one, it's easy to get that

$$(1 - \delta_s) \| x_k - x^* \|_2^2 \le \| A_n(x_k - x^*) \|_2^2 \le (1 + \delta_s) \| x_k - x^* \|_2^2,$$
(11)

and  $\kappa \leq \frac{1+\delta_s}{1-\delta_s} \doteq \kappa_0$ . Substituting this into (10), we derive that



$$\| x_k - x^* \|_2 \le 2\sqrt{\kappa_0} \left( \frac{\sqrt{\kappa_0} - 1}{\sqrt{\kappa_0} + 1} \right)^k \| x_0 - x^* \|_2.$$
(12)

## Convergence analysis

Assume that Algorithm X is a convergent iterative scheme which can be used to solve HTP 2. Just like the FHTP, we propose a class of schemes. In Algorithm 4, Algorithm X can be Algorithms 2, 3 or other algorithms.

#### Algorithm 4 Fast Hard Thresholding Pursuit via Algorithm X

```
Require: parameters needed in Algorithm X and l > 0.
```

Initialization:  $x_0 = 0$ for n = 0, 1, 2, ...  $S_{n+1} = C_s(x_n + A^*(y - Ax_n))$   $x_{n+1}^0 = (x_n + A^*(y - Ax_n))_{S_{n+1}}$ Apply Algorithm X to HTP<sub>2</sub> to get  $x_{n+1}^1, ..., x_{n+1}^l$   $x_{n+1} = x_{n+1}^l$ . end for

But no matter which one, Algorithm 4 is convergent provided Algorithm X is convergent.

**Theorem 3.1** Suppose that the sequence  $x_{n+1}^i$  generated by Algorithm X satisfies  $\|x_n^i - x_n^*\| \le \gamma(i) \|x_n^0 - x_n^*\|$ . Then, for any s-sparse x which obeys y = Ax, the sequence  $x_n$  generated by Algorithm 4 satisfies

$$\| x_n - x \|_2 \le \rho^n \| x_0 - x \|_2, \tag{13}$$

 $\begin{array}{ll} \rho = (1+\sqrt{2}\delta_{3s})(\delta_{2s}+\gamma(l))\sqrt{\frac{1+\delta_{2s}}{1-\delta_{2s}}} + (1+\sqrt{2})\delta_{3s}\gamma(l) + \sqrt{2}\delta_{3s} \\ \text{Because Algorithm $X$ is convergent, } \lim_{l\to\infty}\gamma(l) = 0 \\ \text{.} \\ \text{Let } l\to\infty \end{array}$ 

Because Algorithm X is convergent,  $\lim_{l \to \infty} \gamma(l) = 0$ . Let  $l \to \infty$ ,  $\rho^0 = (1+\sqrt{2}\delta_{3s})\delta_{2s}\sqrt{\frac{1+\delta_{2s}}{1-\delta_{2s}}} + \sqrt{2}\delta_{3s}$ . Note that  $0 < \delta_{2s} < \delta_{3s} < 1$ , it's easy to get  $\rho^0 < (1+\sqrt{2}\delta_{3s})\delta_{3s}\sqrt{\frac{1+\delta_{3s}}{1-\delta_{3s}}} + \sqrt{2}\delta_{3s}$ . If  $\delta_{3s} < 0.3014$ , then  $\rho^0 < 1$ . The numerical experiments demonstrate Algorithm 4 can work as well as FHTP in reconstruction ability.

#### **Numerical Results**

In this section, we report numerical results about Algorithm 4. In the experiments, 'X' will be substituted by the Nesterov method and CG method. All the experiments are run in MatLab 2010a on a server with 2.60GHz Intel processor and 3.90 GB RAM.

Table 1: Time cost by different algorithms. The sensing matrix  $A \in \mathcal{R}^{2000 \times 10000}$ 

scenario	BP	CoSaMP	HTP	FHTP	X=Nesterov	X=CG
Gaussian ma	trices and Gaussia	n vectors				
s=100	690.6s	0.7s	15.5s	15.3s	15.2s	15.3s
s=200	937.2s	11.4s	15.4s	15.2s	15.1s	15.2s
s=300	992.3s	51.8s	15.6s	15.5s	15.3s	15.5s
s=400	1385.8s	136.8s	16.1s	15.8s	15.4s	15.7s
Gaussian ma	trices and Bernou	lli vectors				
s=100	409.6s	0.2s	13.2s	13.9s	11.8s	12.2s
s=200	444.5s	3.2s	15.4s	15.2s	15.1s	14.7s
s=300	447.2s	12.9s	13.7s	14.6s	12.8s	13.1s
s=400	453.7s	32.2s	14.5s	16.9s	13.7s	13.9s
Bernoulli ma	trices and Gaussia	an vectors				
s=100	645.5s	2.6s	14.9s	19.6s	16.5s	19.9s
s=200	638.2s	52.7s	13.5	15.5	13.7s	15.2s
s=300	784.8s	176.9s	13.6s	19.5s	14.5s	16.6s
s=400	1314.1s	1507.5s	14.3s	24.3s	15.4s	19.3s
Bernoulli ma	trices and Bernou	lli vectors				
s=100	330.3s	0.6s	13.5s	15.2s	13.6s	13.8s
s=200	351.1s	7.5s	12.9s	16.8s	14.0s	16.3s
s=300	419.2s	665.2s	13.2s	19.1s	15.8s	18.4s
s=400	406.1s	1531.9s	13.9s	21.1s	18.7s	19.8s



Table 2: Time cost by different algorithms. The sensing matrix  $A \in \mathcal{R}^{4000 \times 10000}$ 

scenario	BP	CoSaMP	HTP	FHTP	X=Nesterov	X=CG
	rices and Gaussia		1111	11111	A-INCSICIOV	A-CG
Gaussian mai	ilices alid Gaussia	ii vectors		•		
s=100	992.9s	1.1s	21.6s	21.4s	20.9s	21.5s
s=200	1117.2s	10.1s	22.0s	22.0s	21.8s	22.1s
s=300	1174.7s	30.9s	22.4s	22.1s	22.0s	21.7s
s=400	1226.3s	34.1s	22.9s	23.4s	22.1s	22.1s
Gaussian mat	rices and Bernoul	li vectors				
s=100	760.5s	0.2s	21.1s	21.5s	21.3s	21.5s
s=200	766.9s	9.3s	21.2s	21.9s	21.4s	21.6s
s=300	760.3.2s	30.6s	22.0s	22.8s	21.3s	21.4s
s=400	765.5s	32.2s	14.5s	16.9s	13.7s	13.9s
Bernoulli mat	trices and Gaussia	n vectors				
s=100	944.2s	2.1s	21.9s	25.1s	22.1s	25.2s
s=200	1116.9s	21.2	22.5s	27.5s	23.0s	26.4s
s=300	1175.7s	104.0s	22.7s	25.9s	23.7s	27.2s
s=400	1220.1s	369.1s	27.3s	31.8s	24.2s	28.5s
Bernoulli mat	trices and Bernoul	li vectors				
s=100	643.4s	0.7s	21.9s	25.6s	22.7s	24.8s
s=200	705.7s	8.3	22.6s	25.8s	22.3s	25.2s
s=300	775.4s	21.3s	22.3s	24.5s	23.0s	27.7s
s=400	856.1s	33.2s	25.2s	31.6s	24.0s	31.4s

Table 3: Time cost by different algorithms. The sensing matrix  $A \in \mathcal{R}^{2000 \times 12000}$ 

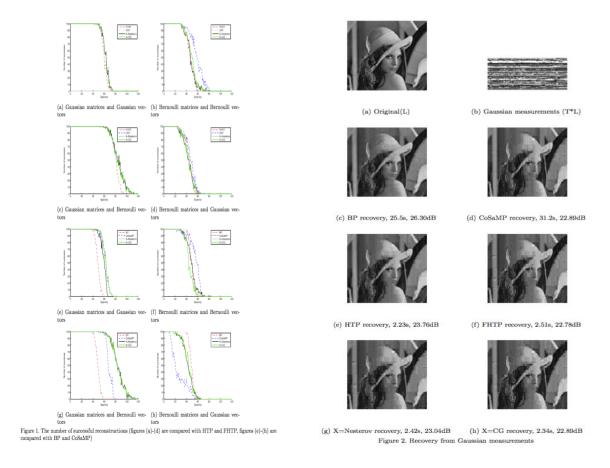
scenario	BP	CoSaMP	HTP	FHTP	X=Nesterov	X=CG
	trices and Gaussia				1	
s=100	897.4s	1.1s	18.4s	17.0s	16.4s	16.6s
s=200	1099.7s	11.2s	17.1s	17.5s	16.8s	17.0s
s=300	1302.8s	42.6s	17.9s	18.4s	17.6s	18.1s
s=400	1506.2s	138.8s	18.1s	18.2s	17.7s	18.1s
Gaussian mat	trices and Bernoul	li vectors				
s=100	459.1s	0.2s	16.3s	16.8s	16.4s	16.6s
s=200	520.4s	3.3s	16.1s	17.3s	16.4s	16.6s
s=300	610.7s	13.0s	16.7s	23.8s	17.1s	17.8s
s=400	680.0s	30.1s	16.9s	21.9s	18.3s	18.4s
Bernoulli ma	trices and Gaussia	n vectors				
s=100	877.9s	5.1s	16.1s	20.2s	17.6s	20.2s
s=200	1045.0s	254.3	17.8s	22.7s	19.0s	22.3s
s=300	1414.1s	994.1s	18.2s	24.9s	22.3s	26.4s
s=400	1773.9s	1722.8s	18.7s	28.7s	21.7s	25.3s
Bernoulli ma	trices and Bernoul	li vectors				
s=100	472.6s	0.7s	16.7s	21.7s	17.7s	19.9s
s=200	468.7s	9.1	16.7s	23.6s	19.2s	22.1s
s=300	656.8s	727.0s	17.6s	24.6s	22.2s	24.7s
s=400	1500.0s	1363.0s	19.7s	32.3s	24.3s	31.2s

I) The first experiment is about recovery ability of Algorithm 4. To show the results clearly, this experiment contains two parts: one focuses on comparing Algorithm 4 with FHTP and HTP; the other one focuses on comparing Algorithm 4 with other classical algorithms, namely the BP, CoSaMP algorithms (the matlab codes of these algorithms can be found in [18]). In this experiment, the sensing matrix A is a  $200 \times 1000$  Gaussian matrix or Bernoulli matrix and S-sparse Gaussian or Bernoulli signals are employed. We run the algorithms with a stopping criterion  $\|y-Ax_n\|_2 \le 10^{-4} \|y\|_2$  or 100 steps. And  $\|x-x_n\|_2 \le 10^{-6} \|x\|_2$  means a successful reconstruction. For each  $s \in [1,120]$ , we use the 100 different matrices and sparse vectors. Figure 1 presents the numerical results. The results of experiment 1 show that the Algorithm 4 have proved capable of recovering the sparse signal from a small amount of linear observations. Just as what we declared before, the Algorithm 4 can work well in numerical tests.

2) The second experiment is about the time cost by Algorithm 4. In this test, we use three sensing matrices: 2000×10000, 4000×10000, and 2000×12000. In this test, we concern on comparing



Algorithm 4 with the BP, CoSaMP, HTP and FHTP algorithm. The stopping criterion can be only set as  $\|y-Ax_n\|_2 \le 10^{-4} \|y\|_2$  for that the reconstructions are all successful. Tables 1, 2 and 3 present the numerical results. Form the results of experiment 2, we can find that with almost the same accuracy Algorithm 4 with "X=Nesterov" and "X=CG" are faster than FHTP and BP. When the sparsity decreases, Algorithm 4 is also faster than CoSaMP. As the scale of the problem increases, more time can be saved.



3) The third experiment is conducted on a real-world example. We employ the *Lena* picture  $(L \in \mathcal{R}^{256 \times 256})$  and the Haar wavelet frame matrix H[17] for test. We concern on reconstructing L from linear measurements G = TL, where  $T \in \mathcal{R}^{120 \times 256}$  is a Gaussian matrix. As is well known, L can be expressed as L = HX, where  $X \in \mathcal{R}^{256 \times 256}$  and most entries of X are relative small. Therefore, X can be regarded as a sparse matrix. Then, the problem reduces to recovering sparse X from G = (TH)X.

The BP, CoSaMP, HTP, FHTP and Algorithm 4 are used to reconstruct the picture. The parameter S is set as S = 20. Figure 2 presents the reconstructed picture, the time cost, and the Signal-to-Noise Ratio(SNR).

#### Conclusion

In this paper, we propose a class of schemes for sparse recovery. The FHTP method can be regarded as a special one of this class. Theoretically, these proposed algorithms are proved to converge linearly. Finally, the numerical result validates the feasibility and efficiency of the proposed schemes.



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